

From Algebraic Geometry to Linear Cryptanalysis

Application to Anemoi

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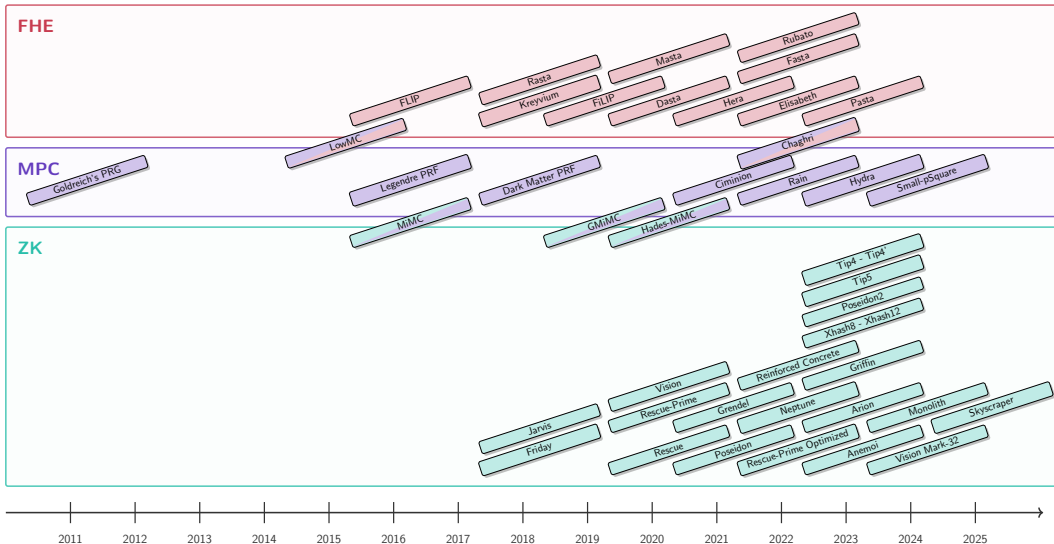
(joint work with Tim Beyne)



ALPSY Workshop, Obergurgl, Austria
January 26th, 2025



New symmetric primitives



A new context

Traditional case

Alphabet

Operations based on logical gates or CPU instructions.

$$\mathbb{F}_2^n, \text{ with } n \simeq 4, 8$$

Arithmetization-Oriented

Alphabet

Operations based on large finite-field arithmetic.

$$\mathbb{F}_q, \text{ with } q \in \{2^n, p\}, p \simeq 2^n, n \geq 32$$

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Cryptanalysis

Decades of cryptanalysis

- ★ algebraic attacks ✓
- ★ differential attacks ✓
- ★ linear attacks ✓
- ★ ...

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Cryptanalysis

≤ 8 years of cryptanalysis

- ★ algebraic attacks ✓
- ★ differential attacks ✗
- ★ linear attacks ✗
- ★ ...

Linearity

Definition

Let $F : \mathbb{F}_q^n \rightarrow \mathbb{F}_q^m$ be a function and ω a primitive character. The **Walsh transform** for the character ω of the linear approximation (u, v) of F is given by

$$\mathcal{W}_{u,v}^F = \sum_{x \in \mathbb{F}_q^n} \omega(\langle v, F(x) \rangle - \langle u, x \rangle) .$$

$$\mathcal{W}_{u,v}^F = q^n \cdot C_{u,v}^F$$

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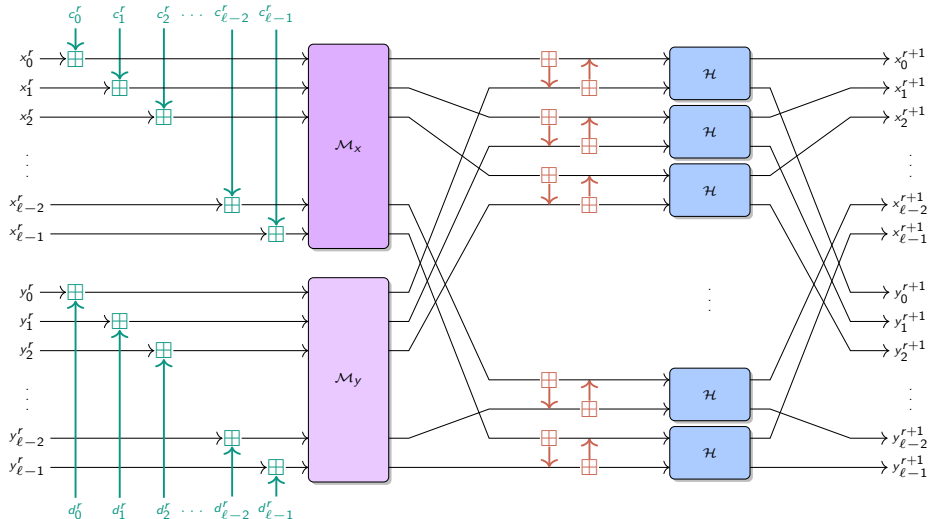
Definition

The **Linearity** \mathcal{L}_F of $F : \mathbb{F}_q^n \rightarrow \mathbb{F}_q^m$ is the highest Walsh coefficient.

$$\mathcal{L}_F = \max_{u,v \neq 0} |\mathcal{W}_{u,v}^F| .$$

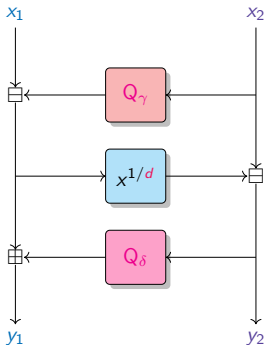
Anemoi round function

Introduced by [Bouvier, Briaud, Chaidos, Perrin, Salen, Velichkov and Willems, 2023].



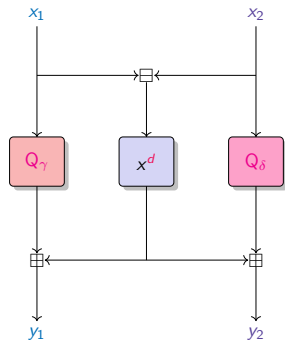
Flystel - Definition

Introduced by [Bouvier, Briaud, Chaidos, Perrin, Salen, Velichkov and Willems, 2023].



Open variant.

$$\begin{cases} y_1 = x_1 - Q_\gamma(x_2) + Q_\delta(x_2 - (x_1 - Q_\gamma(x_2))^{1/d}) \\ y_2 = x_2 - (x_1 - Q_\gamma(x_2))^{1/d}. \end{cases}$$

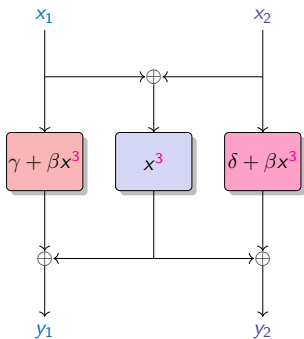


Closed variant.

$$\begin{cases} y_1 = (x_1 - x_2)^d + Q_\gamma(x_1) \\ y_2 = (x_1 - x_2)^d + Q_\delta(x_2). \end{cases}$$

Closed Flystel in \mathbb{F}_{2^n}

Introduced by [Bouvier, Briaud, Chaidos, Perrin, Salen, Velichkov and Willems, 2023].



Closed Flystel.

$$\mathcal{L}_F = \max_{u, v \neq 0} \left| \sum_{x \in \mathbb{F}_{2^n}} (-1)^{(\langle v, F(x) \rangle - \langle u, x \rangle)} \right|$$

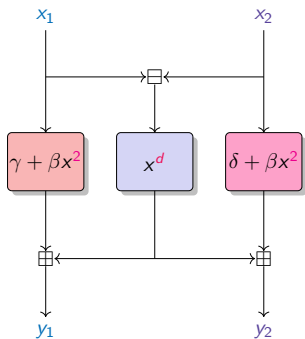
Bound

Linearity bound for the Flystel:

$$\mathcal{L}_F \leq 2^{n+1}$$

Closed Flystel in \mathbb{F}_p

Introduced by [Bouvier, Briaud, Chaidos, Perrin, Salen, Velichkov and Willems, 2023].



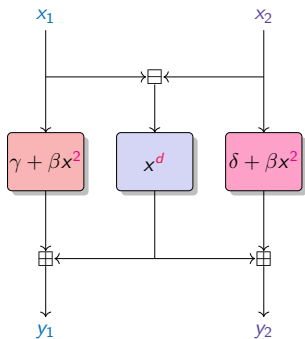
Closed Flystel.

d is a small integer s.t.
 $x \mapsto x^d$ is a permutation of \mathbb{F}_p
 (usually $d = 3, 5$).

$$\mathcal{L}_F = \max_{u, v \neq 0} \left| \sum_{x \in \mathbb{F}_p^2} e\left(\frac{2i\pi}{p}\right) (\langle v, F(x) \rangle - \langle u, x \rangle) \right|$$

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How to determine an accurate bound for the linearity of the Closed Flystel in \mathbb{F}_p ?

Weil bound

Proposition [Weil, 1948]

Let $f \in \mathbb{F}_p[x]$ be a univariate polynomial with $\deg(f) = d$. Then

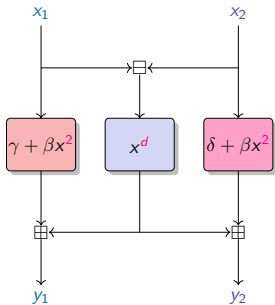
$$\mathcal{L}_f \leq (d - 1)\sqrt{p}$$

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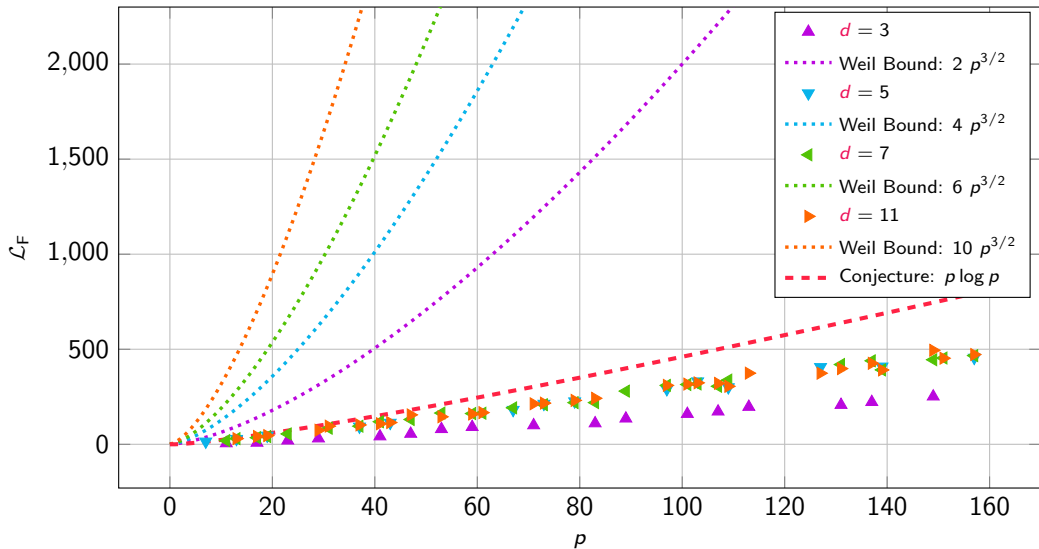
Closed Flystel.

$$\mathcal{L}_F \leq (d - 1)p\sqrt{p} ? \quad \begin{cases} \mathcal{L}_{\gamma+\beta x^2} \leq \sqrt{p}, \\ \mathcal{L}_{x^d} \leq (d - 1)\sqrt{p}, \\ \mathcal{L}_{\delta+\beta x^2} \leq \sqrt{p}. \end{cases}$$

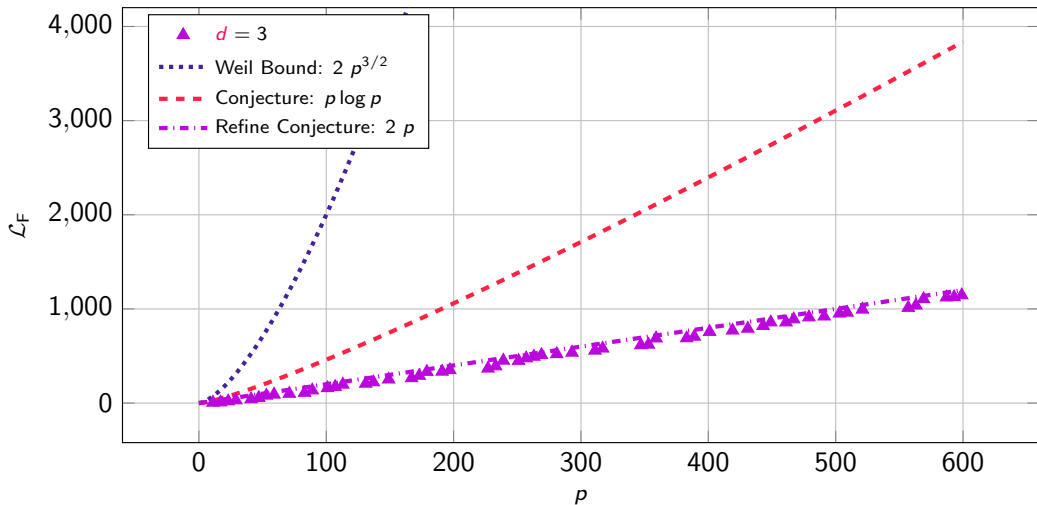
Conjecture

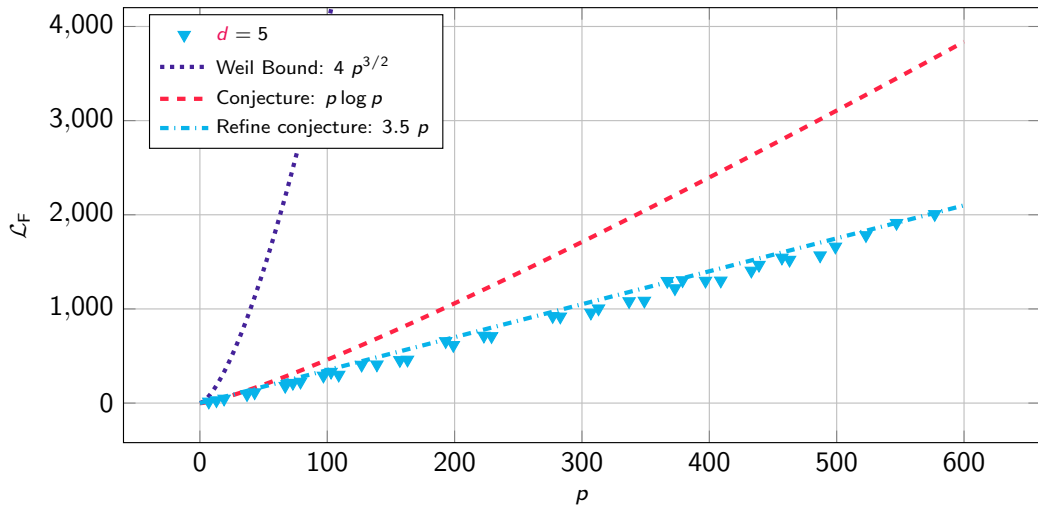
$$\mathcal{L}_F = \max_{u, v \neq 0} \left| \sum_{x \in \mathbb{F}_p^2} e\left(\frac{2i\pi}{p}\right) (\langle v, F(x) \rangle - \langle u, x \rangle) \right| \leq p \log p$$

Experimental results



Experimental results ($d = 3$)



Experimental results ($d = 5$)

Take-away

AO primitives: new symmetric primitives defined over prime fields.

Need for new linear cryptanalysis tools

Take-away

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This Talk:

- ★ Applications of results for exponential sums (generalization of Weil bound)

$$\mathcal{W}_{u,v}^F = \sum_{x \in \mathbb{F}_q^n} \omega(\langle v, F(x) \rangle - \langle u, x \rangle) \quad \rightarrow \quad S(f) = \sum_{x \in \mathbb{F}_q^n} \omega^{f(x)} .$$

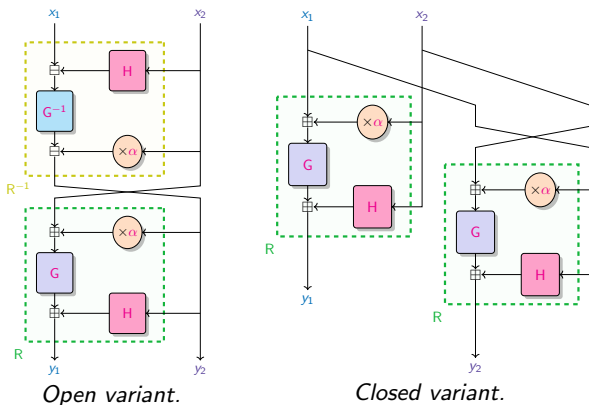
- ★ \mathbb{F}_q is a finite field s.t. q is a power of a prime p .
- ★ Functions with 2 variables $F \in \mathbb{F}_q[x_1, x_2]$.

Generalizations of Weil bound

- ★ **Deligne** bound
 - ★ Application to the **Generalized Butterfly** construction
- ★ **Denef and Loeser** bound
 - ★ Application to **3-round Feistel** construction
- ★ **Rojas-León** bound
 - ★ Application to the **Generalized Flystel** construction

Generalized Butterfly

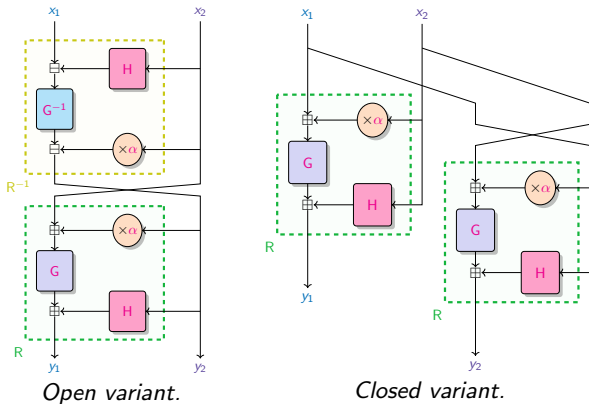
Originally introduced by [Perrin, Udovenko and Biryukov, 2016] over binary fields, \mathbb{F}_{2^n} , n odd.
 BUTTERFLY[G, H, α], with $G : \mathbb{F}_q \rightarrow \mathbb{F}_q$ a permutation, $H : \mathbb{F}_q \rightarrow \mathbb{F}_q$ a function and $\alpha \in \mathbb{F}_q$.



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$$f(x_1, x_2) = \langle v, F(x) \rangle - \langle u, x \rangle = v_1 G(x_1 + \alpha x_2) + v_2 G(x_2 + \alpha x_1) + v_1 H(x_2) + v_2 H(x_1) - u_1 x_1 - u_2 x_2 .$$



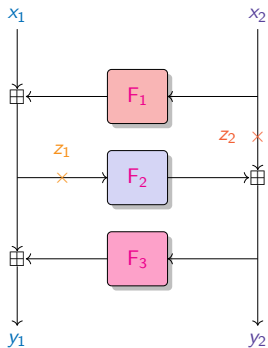
Deligne Bound

The hypersurface defined by $f_d = 0$ (the degree- d homogeneous component of f) is smooth so that

$$\mathcal{L}_F \leq (\max\{\deg G, \deg H\} - 1)^2 \cdot q$$

3-round Feistel

Let $\text{FEISTEL}[F_1, F_2, F_3]$ be a 3-round Feistel network with F_2 a permutation and $d_1 \geq d_3$ where $d_1 = \deg(F_1)$, $d_2 = \deg(F_2)$, and $d_3 = \deg(F_3)$.

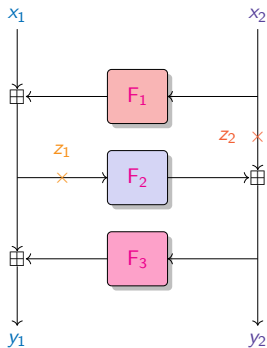


A 3-round Feistel.

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$$f(z_1, z_2) = \langle v, F(z) \rangle - \langle u, z \rangle = v_1 F_3(z_2 + F_2(z_1)) + v_2 F_2(z_1) + u_1 F_1(z_2) + (v_1 - u_1)z_1 + (v_2 - u_2)z_2.$$

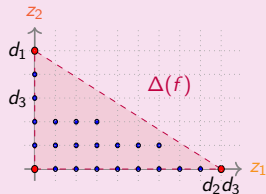


A 3-round Feistel.

Denef-Loeser Bound

f is commode and non-degenerate, with Newton number:

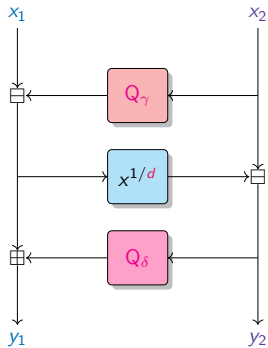
$$\nu(f) = (d_1 - 1)(d_2 d_3 - 1).$$



$$\mathcal{L}_F \leq (d_1 - 1)(d_2 d_3 - 1) \cdot q$$

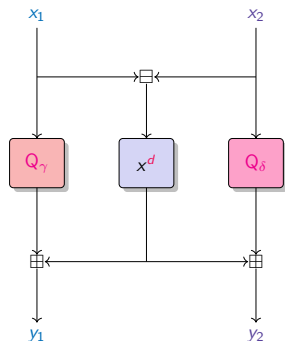
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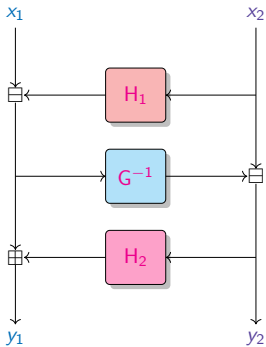


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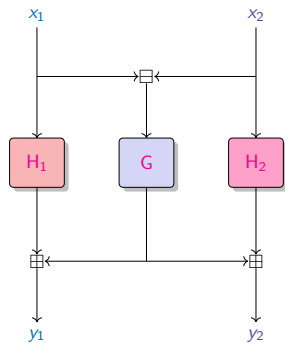
$$\begin{cases} y_1 &= (x_1 - x_2)^d + Q_\gamma(x_1) \\ y_2 &= (x_1 - x_2)^d + Q_\delta(x_2). \end{cases}$$

Generalized Flystel - Definition

$F = \text{FLYSTEL}[H_1, G, H_2]$, with $G : \mathbb{F}_q \rightarrow \mathbb{F}_q$ a permutation, and $H_1, H_2 : \mathbb{F}_q \rightarrow \mathbb{F}_q$ functions.



Open variant.



Closed variant.

$$\begin{cases} y_1 = x_1 - H_1(x_2) + H_2(x_2 - G^{-1}(x_1 - H_1(x_2))) \\ y_2 = x_2 - G^{-1}(x_1 - H_1(x_2)). \end{cases}$$

$$\begin{cases} y_1 = G(x_1 - x_2) + H_1(x_1) \\ y_2 = G(x_1 - x_2) + H_2(x_2). \end{cases}$$

Isolated singularities

Definition

- ★ A singular point of a hypersurface is **isolated** if there exists a Zariski neighborhood of the point that contains no other singular points.
- ★ A polynomial g is **quasi-homogeneous** of degree δ if there exists w_1, \dots, w_n s.t.

$$g(\lambda^{w_1} x_1, \dots, \lambda^{w_n} x_n) = \lambda^\delta g(x_1, \dots, x_n) .$$

- ★ The **Milnor number** of the singularity is equal to $\prod_{i=1}^n (\delta/w_i - 1)$

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Example: Let $f(x) = (x - 1)^d$.

- ★ $x = 1$ is the **only singular point** of $f = 0$.
- ★ Up to translation, we can consider the singularity in the origin: $g(x) = x^d$.

$$g(\lambda^w x) = (\lambda^w x)^d = \lambda^{w \cdot d} x^d = \lambda^{w \cdot d} g(x) \quad \text{so that } \delta = w \cdot d$$

- ★ **Milnor number** of the singularity: $\delta/w - 1 = d - 1$.

Rojas-León Theorem

Theorem [Rojas-León, 2006]

Let $f \in \mathbb{F}_q[x_1, \dots, x_n]$, s.t. $\deg(f) = d$.

Suppose that $f = f_d + f_{d'} + \dots$, where $f_d, f_{d'}$, are resp. **the degree- d , degree- d' , homogeneous component** of f , with $\gcd(d, p) = \gcd(d', p) = 1$ and $d'/d > p/(p + (p - 1)^2)$.

If the following conditions are satisfied

- ★ the hypersurface defined by $f_d = 0$ has at worst **quasi-homogeneous isolated singularities** of degrees prime to p with **Milnor numbers** μ_1, \dots, μ_s ,
- ★ the hypersurface defined by $f_{d'} = 0$ contains none of these singularities,

then we have

$$|S(f)| = \left| \sum_{x \in \mathbb{F}_q^n} \omega^{f(x)} \right| \leq \left((d-1)^n - (d-d') \sum_{i=1}^s \mu_i \right) \cdot q^{n/2}.$$

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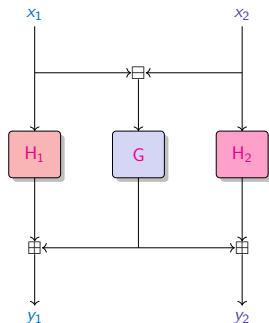
$$|S(f)| = \left| \sum_{x \in \mathbb{F}_q^n} \omega^{f(x)} \right| \leq \left((d-1)^n - (d-d') \sum_{i=1}^s \mu_i \right) \cdot q^{n/2}.$$

$$\text{Linearity bound for } n = 2: \mathcal{L}_F \leq ((d-1)^2 - (d-d') \sum_{i=1}^s \mu_i) \cdot q.$$

Generalized Flystel - Bound

Let $F = \text{FLYSTEL}[H_1, G, H_2]$, with G a permutation, H_1, H_2 functions ($\deg G > \deg H_1, \deg H_2$).

$$\begin{aligned} f(x_1, x_2) &= \langle (v_1, v_2), F(x_1, x_2) \rangle - \langle (u_1, u_2), (x_1, x_2) \rangle \\ &= (v_1 + v_2) G(x_1 - x_2) + v_1 H_1(x_1) + v_2 H_2(x_2) - u_1 x_1 - u_2 x_2 . \end{aligned}$$

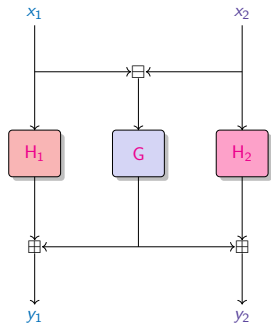


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Linearity Bound

- ★ The hypersurface

$$f_d = (v_1 + v_2)(x_1 - x_2)^d = 0$$

contains one singular point $[1 : 1]$ of quasi-homogeneous type with Milnor number $d - 1$.

- ★ The hypersurface

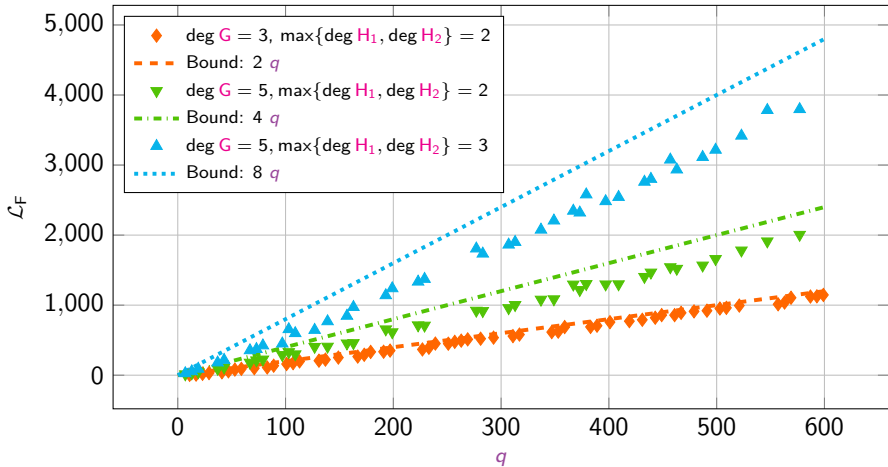
$$f_{d'} = v_i x_i^{\deg H_i} = 0$$

does not contain this point.

$$\mathcal{L}_F \leq (\deg G - 1)(\max\{\deg H_1, \deg H_2\} - 1) \cdot q$$

Generalized Flystel - Results

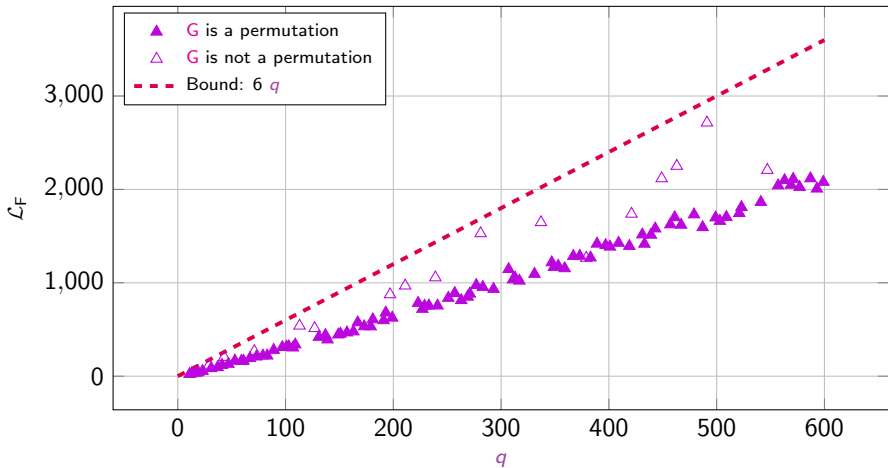
Let $F = \text{FLYSTEL}[H_1, G, H_2]$ with H_1 , G and H_2 monomials.



Low-degree permutations G , H_1 and H_2 .

Generalized Flystel - Results

Let $F = \text{FLYSTEL}[H_1, G, H_2]$ with H_1 , G and H_2 monomials.



$\deg G = 7$ and $\deg H_1 = \deg H_2 = 2$.

Solving conjecture

Conjecture

Let $F = \text{FLYSTEEL}[\mathbf{H}_1, \mathbf{G}, \mathbf{H}_2]$ be defined by $\mathbf{H}_1(x) = \gamma + \beta x^2$, $\mathbf{G}(x) = x^d$ and $\mathbf{H}_2 = \delta + \beta x^2$, with $\gamma, \delta \in \mathbb{F}_p$ and $\beta \in \mathbb{F}_p^\times$. Then

$$\mathcal{L}_F \leq p \log p .$$

Solving conjecture

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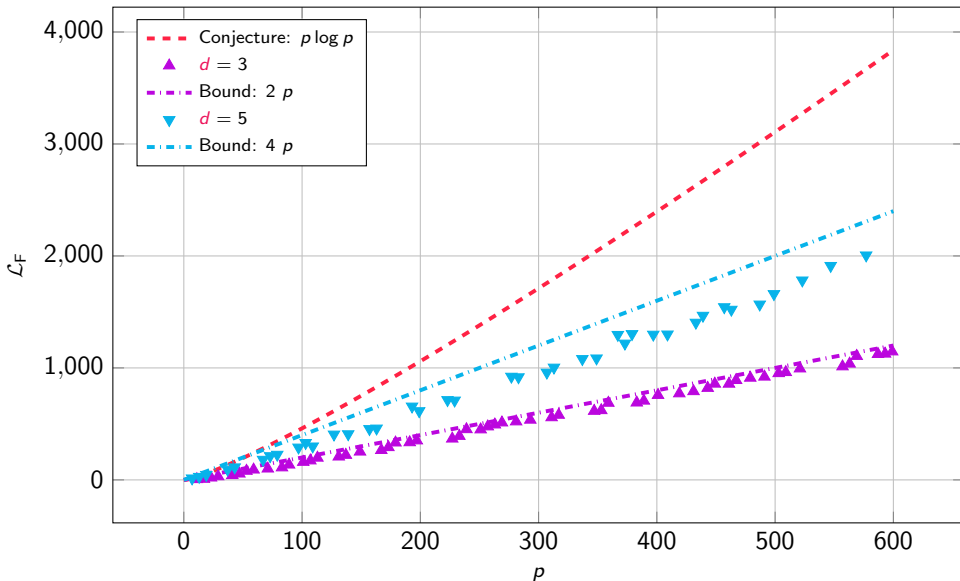
Conjecture proved for $d \leq \log p$

Proposition

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$$\mathcal{L}_F \leq (d - 1)p .$$

Solving conjecture



Cohomological framework

$$S(f) = \sum_{x \in \mathbb{F}_q^n} \omega^{f(x)} = \sum_{x \in \mathbb{F}_q^n} \omega(\langle v, F(x) \rangle - \langle u, x \rangle)$$

Cohomological framework

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Cohomological framework



$$|S(f)| = \left| \sum_{i=0}^{2n} (-1)^i \text{Tr}(F | H_c^i(\mathbb{A}^n, \mathcal{L})) \right|$$

Sum of **traces** of the **Frobenius automorphism** on ℓ -adic cohomology groups.

Cohomological framework

$$S(f) = \sum_{x \in \mathbb{F}_q^n} \omega^{f(x)} = \sum_{x \in \mathbb{F}_q^n} \omega(\langle v, F(x) \rangle - \langle u, x \rangle)$$



Cohomological framework



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Thank you

